

MIPS Assembly Language Programming

COE 308

Computer Architecture

Prof. Muhamed Mudawar

College of Computer Sciences and Engineering
King Fahd University of Petroleum and Minerals

Presentation Outline

- ❖ **Assembly Language Statements**
- ❖ Assembly Language Program Template
- ❖ Defining Data
- ❖ Memory Alignment and Byte Ordering
- ❖ System Calls
- ❖ Procedures
- ❖ Parameter Passing and the Runtime Stack

Assembly Language Statements

- ❖ Three types of statements in assembly language
 - ❖ Typically, one statement should appear on a line
- 1. Executable Instructions
 - ❖ Generate machine code for the processor to execute at runtime
 - ❖ Instructions tell the processor what to do
- 2. Pseudo-Instructions and Macros
 - ❖ Translated by the assembler into real instructions
 - ❖ Simplify the programmer task
- 3. Assembler Directives
 - ❖ Provide information to the assembler while translating a program
 - ❖ Used to define segments, allocate memory variables, etc.
 - ❖ Non-executable: directives are not part of the instruction set

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Instructions

- ❖ Assembly language instructions have the format:
`[label:] mnemonic [operands] [#comment]`
- ❖ Label: (optional)
 - ❖ Marks the address of a memory location, must have a colon
 - ❖ Typically appear in data and text segments
- ❖ Mnemonic
 - ❖ Identifies the operation (e.g. `add`, `sub`, etc.)
- ❖ Operands
 - ❖ Specify the data required by the operation
 - ❖ Operands can be registers, memory variables, or constants
 - ❖ Most instructions have three operands

```
L1:    addiu $t0, $t0, 1           #increment $t0
```

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Comments

- ❖ Comments are very important!
 - ❖ Explain the program's purpose
 - ❖ When it was written, revised, and by whom
 - ❖ Explain data used in the program, input, and output
 - ❖ Explain instruction sequences and algorithms used
 - ❖ Comments are also required at the beginning of every procedure
 - Indicate input parameters and results of a procedure
 - Describe what the procedure does
- ❖ Single-line comment
 - ❖ Begins with a hash symbol **#** and terminates at end of line

Next ...

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Program Template

```
# Title:                               Filename:
# Author:                              Date:
# Description:
# Input:
# Output:
##### Data segment #####
.data
. . .
##### Code segment #####
.text
.globl main
main:                                  # main program entry
. . .
li $v0, 10                             # Exit program
syscall
```

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.DATA, .TEXT, & .GLOBL Directives

❖ .DATA directive

- ✧ Defines the **data segment** of a program containing data
- ✧ The program's variables should be defined under this directive
- ✧ Assembler will allocate and initialize the storage of variables

❖ .TEXT directive

- ✧ Defines the **code segment** of a program containing instructions

❖ .GLOBL directive

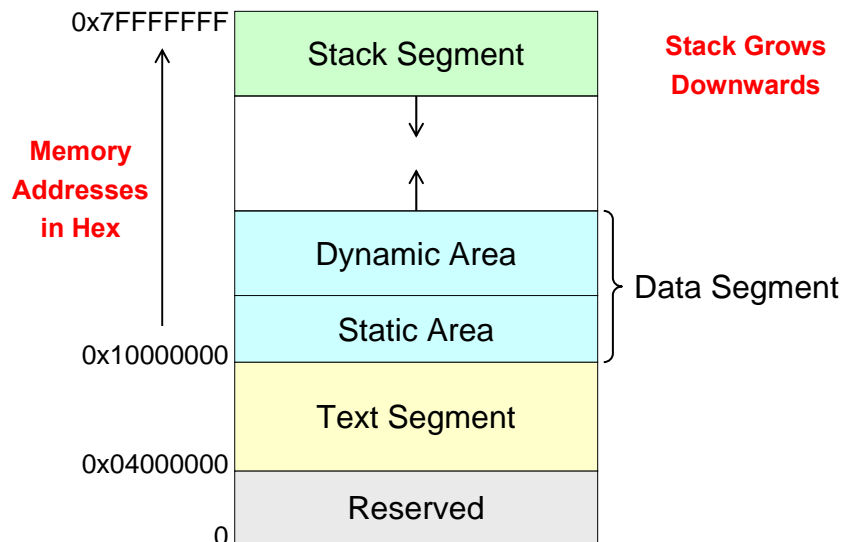
- ✧ Declares a symbol as **global**
- ✧ Global symbols can be referenced from other files
- ✧ We use this directive to declare *main* procedure of a program

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Layout of a Program in Memory



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Data Definition Statement

- ❖ Sets aside storage in memory for a variable
- ❖ May optionally assign a name (label) to the data
- ❖ Syntax:

[*name*:] *directive* *initializer* [, *initializer*] . . .



var1: **.WORD** **10**

- ❖ All initializers become binary data in memory

Data Directives

- ❖ **.BYTE** Directive
 - ❖ Stores the list of values as 8-bit bytes
- ❖ **.HALF** Directive
 - ❖ Stores the list as 16-bit values aligned on half-word boundary
- ❖ **.WORD** Directive
 - ❖ Stores the list as 32-bit values aligned on a word boundary
- ❖ **.FLOAT** Directive
 - ❖ Stores the listed values as single-precision floating point
- ❖ **.DOUBLE** Directive
 - ❖ Stores the listed values as double-precision floating point

String Directives

❖ **.ASCII** Directive

- ❖ Allocates a sequence of bytes for an ASCII string

❖ **.ASCIIZ** Directive

- ❖ Same as **.ASCII** directive, but adds a NULL char at end of string
- ❖ Strings are null-terminated, as in the C programming language

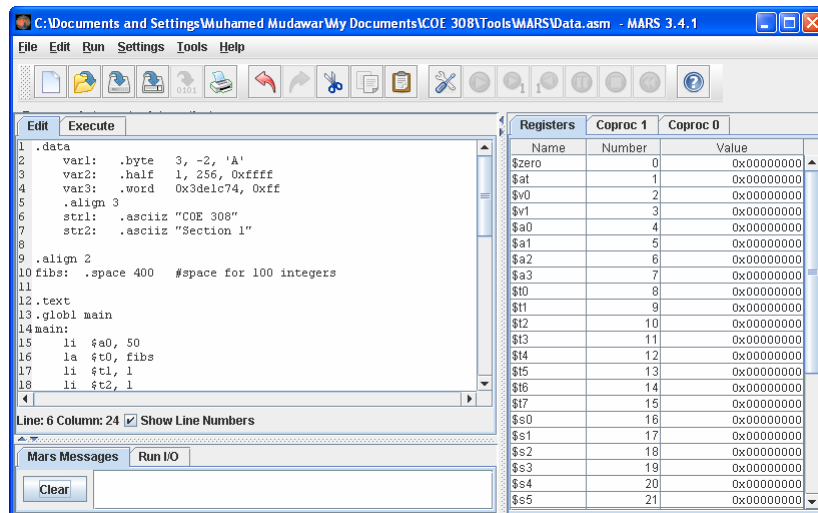
❖ **.SPACE** Directive

- ❖ Allocates space of n uninitialized bytes in the data segment

Examples of Data Definitions

```
.DATA
var1:  .BYTE    'A', 'E', 127, -1, '\n'
var2:  .HALF    -10, 0xffff
var3:  .WORD    0x12345678
var4:  .FLOAT   12.3, -0.1
var5:  .DOUBLE  1.5e-10
str1:  .ASCII   "A String\n"
str2:  .ASCIIZ  "NULL Terminated String"
array: .SPACE   100
```

MARS Assembler and Simulator Tool

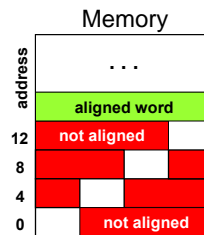


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Memory Alignment

- ❖ Memory is viewed as an **array of bytes** with addresses
 - ❖ **Byte Addressing**: address points to a byte in memory
- ❖ Words occupy 4 consecutive bytes in memory
 - ❖ MIPS instructions and integers occupy 4 bytes
- ❖ **Alignment: address is a multiple of size**
 - ❖ Word address should be a multiple of **4**
 - Least significant 2 bits of address should be **00**
 - ❖ Halfword address should be a multiple of **2**
- ❖ **.ALIGN n** directive
 - ❖ Aligns the next data definition on a 2^n byte boundary



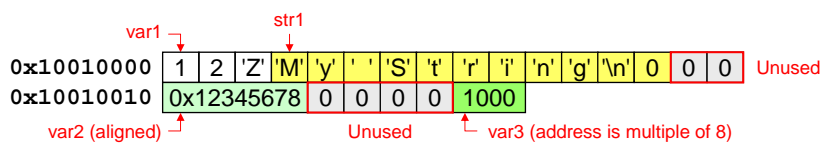
Symbol Table

- ❖ Assembler builds a **symbol table** for labels (variables)
 - ❖ Assembler computes the address of each label in data segment
- ❖ Example

Symbol Table

```
.DATA
var1: .BYTE    1, 2, 'Z'
str1: .ASCIIZ  "My String\n"
var2: .WORD    0x12345678
.ALIGN 3
var3: .HALF    1000
```

Label	Address
var1	0x10010000
str1	0x10010003
var2	0x10010010
var3	0x10010018



System Calls

- ❖ Programs do input/output through system calls
- ❖ MIPS provides a special **syscall** instruction
 - ❖ To obtain services from the operating system
 - ❖ Many services are provided in the SPIM and MARS simulators
- ❖ Using the **syscall** system services
 - ❖ Load the service number in register **\$v0**
 - ❖ Load argument values, if any, in registers **\$a0**, **\$a1**, etc.
 - ❖ Issue the **syscall** instruction
 - ❖ Retrieve return values, if any, from result registers

Syscall Services

Service	\$v0	Arguments / Result
Print Integer	1	\$a0 = integer value to print
Print Float	2	\$f12 = float value to print
Print Double	3	\$f12 = double value to print
Print String	4	\$a0 = address of null-terminated string
Read Integer	5	\$v0 = integer read
Read Float	6	\$f0 = float read
Read Double	7	\$f0 = double read
Read String	8	\$a0 = address of input buffer \$a1 = maximum number of characters to read
Exit Program	10	
Print Char	11	\$a0 = character to print
Read Char	12	\$a0 = character read

Supported by MARS

Reading and Printing an Integer

```
##### Code segment #####  
.text  
.globl main  
main:                                # main program entry  
    li    $v0, 5                      # Read integer  
    syscall                            # $v0 = value read  
  
    move  $a0, $v0                    # $a0 = value to print  
    li    $v0, 1                      # Print integer  
    syscall  
  
    li    $v0, 10                     # Exit program  
    syscall
```

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Reading and Printing a String

```
##### Data segment #####  
.data  
    str: .space 10                    # array of 10 bytes  
##### Code segment #####  
.text  
.globl main  
main:                                # main program entry  
    la    $a0, str                    # $a0 = address of str  
    li    $a1, 10                     # $a1 = max string length  
    li    $v0, 8                      # read string  
    syscall  
    li    $v0, 4                      # Print string str  
    syscall  
    li    $v0, 10                     # Exit program  
    syscall
```

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Program 1: Sum of Three Integers

```
# Sum of three integers
#
# Objective: Computes the sum of three integers.
# Input: Requests three numbers.
# Output: Outputs the sum.
##### Data segment #####
.data
prompt: .ascii "Please enter three numbers: \n"
sum_msg: .ascii "The sum is: "
##### Code segment #####
.text
.globl main
main:
    la    $a0,prompt        # display prompt string
    li    $v0,4
    syscall
    li    $v0,5              # read 1st integer into $t0
    syscall
    move  $t0,$v0
```

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Sum of Three Integers - Slide 2 of 2

```
    li    $v0,5              # read 2nd integer into $t1
    syscall
    move  $t1,$v0
    li    $v0,5              # read 3rd integer into $t2
    syscall
    move  $t2,$v0
    addu  $t0,$t0,$t1        # accumulate the sum
    addu  $t0,$t0,$t2
    la    $a0,sum_msg        # write sum message
    li    $v0,4
    syscall
    move  $a0,$t0            # output sum
    li    $v0,1
    syscall
    li    $v0,10            # exit
    syscall
```

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Program 2: Case Conversion

```
# Objective: Convert lowercase letters to uppercase
# Input: Requests a character string from the user.
# Output: Prints the input string in uppercase.
##### Data segment #####
.data
name_prompt: .asciiz      "Please type your name: "
out_msg:     .asciiz      "Your name in capitals is: "
in_name:     .space 31    # space for input string
##### Code segment #####
.text
.globl main
main:
    la    $a0,name_prompt # print prompt string
    li    $v0,4
    syscall
    la    $a0,in_name     # read the input string
    li    $a1,31          # at most 30 chars + 1 null char
    li    $v0,8
    syscall
```

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Case Conversion - Slide 2 of 2

```
    la    $a0,out_msg     # write output message
    li    $v0,4
    syscall
    la    $t0,in_name
loop:
    lb    $t1,($t0)
    beqz  $t1,exit_loop   # if NULL, we are done
    blt   $t1,'a',no_change
    bgt   $t1,'z',no_change
    addiu $t1,$t1,-32     # convert to uppercase: 'A'-'a'=-32
    sb    $t1,($t0)
no_change:
    addiu $t0,$t0,1      # increment pointer
    j     loop
exit_loop:
    la    $a0,in_name     # output converted string
    li    $v0,4
    syscall
    li    $v0,10          # exit
    syscall
```

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- ❖ **Procedures**
- ❖ Parameter Passing and the Runtime Stack

Procedures

- ❖ Consider the following `swap` procedure (written in C)
- ❖ Translate this procedure to MIPS assembly language

```
void swap(int v[], int k)
{
    int temp;
    temp = v[k];
    v[k] = v[k+1];
    v[k+1] = temp;
}
```

Parameters:

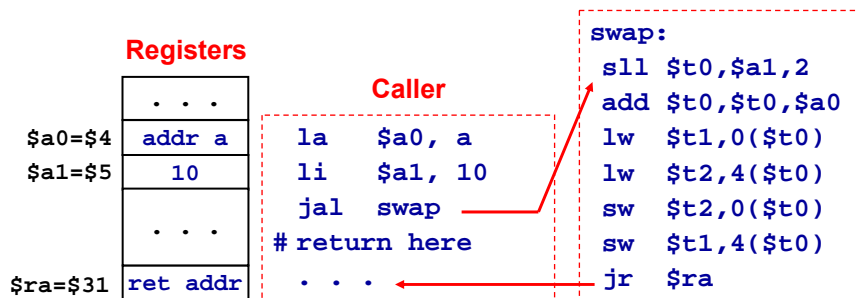
`$a0` = Address of `v[]`
`$a1` = `k`, and
Return address is in `$ra`

swap:

```
sll $t0,$a1,2    # $t0=k*4
add $t0,$t0,$a0  # $t0=v+k*4
lw  $t1,0($t0)   # $t1=v[k]
lw  $t2,4($t0)   # $t2=v[k+1]
sw  $t2,0($t0)   # v[k]=$t2
sw  $t1,4($t0)   # v[k+1]=$t1
jr  $ra          # return
```

Call / Return Sequence

- ❖ Suppose we call procedure swap as: `swap(a, 10)`
 - ❖ Pass **address** of array `a` and `10` as arguments
 - ❖ Call the procedure swap saving **return address** in `$31 = $ra`
 - ❖ Execute procedure swap
 - ❖ Return control to the point of origin (return address)



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Details of JAL and JR

Address	Instructions	Assembly Language	
00400020	lui \$1, 0x1001	la \$a0, a	Pseudo-Direct Addressing PC = imm26 << 2 0x10000f << 2 = 0x0040003C
00400024	ori \$4, \$1, 0		
00400028	ori \$5, \$0, 10	ori \$a1, \$0, 10	
0040002C	jal 0x10000f	jal swap	
00400030	...	# return here	
0040003C	sll \$8, \$5, 2	swap: sll \$t0, \$a1, 2	\$31 = 0x00400030
00400040	add \$8, \$8, \$4	add \$t0, \$t0, \$a0	Register \$31 is the return address register
00400044	lw \$9, 0(\$8)	lw \$t1, 0(\$t0)	
00400048	lw \$10, 4(\$8)	lw \$t2, 4(\$t0)	
0040004C	sw \$10, 0(\$8)	sw \$t2, 0(\$t0)	
00400050	sw \$9, 4(\$8)	sw \$t1, 4(\$t0)	
00400054	jr \$31	jr \$ra	

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Instructions for Procedures

- ❖ JAL (**Jump-and-Link**) used as the call instruction
 - ❖ Save return address in $\$ra = PC+4$ and jump to procedure
 - ❖ Register $\$ra = \31 is used by JAL as the **return address**
- ❖ JR (**Jump Register**) used to return from a procedure
 - ❖ Jump to instruction whose address is in register Rs (PC = Rs)
- ❖ JALR (**Jump-and-Link Register**)
 - ❖ Save return address in Rd = PC+4, and
 - ❖ Jump to procedure whose address is in register Rs (PC = Rs)
 - ❖ Can be used to call methods (addresses known only at runtime)

Instruction	Meaning	Format								
jal	label	$\$31=PC+4$, jump	op ⁶ = 3	imm ²⁶						
jr	Rs	PC = Rs	op ⁶ = 0	rs ⁵	0	0	0	0	8	
jalr	Rd, Rs	Rd=PC+4, PC=Rs	op ⁶ = 0	rs ⁵	0	rd ⁵	0	0	9	

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- ❖ **Parameter Passing and the Runtime Stack**

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Parameter Passing

- ❖ Parameter passing in assembly language is different
 - ❖ More complicated than that used in a high-level language
- ❖ In assembly language
 - ❖ Place all required parameters in an accessible storage area
 - ❖ Then call the procedure
- ❖ Two types of storage areas used
 - ❖ Registers: general-purpose registers are used (**register method**)
 - ❖ Memory: stack is used (**stack method**)
- ❖ Two common mechanisms of parameter passing
 - ❖ Pass-by-value: parameter **value** is passed
 - ❖ Pass-by-reference: **address** of parameter is passed

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Parameter Passing - cont'd

- ❖ By convention, register are used for parameter passing
 - ❖ $\$a0 = \$4 \dots \$a3 = \7 are used for **passing arguments**
 - ❖ $\$v0 = \$2 \dots \$v1 = \3 are used for **result values**
- ❖ Additional arguments/results can be placed on the stack
- ❖ Runtime stack is also needed to ...
 - ❖ Store variables / data structures when they cannot fit in registers
 - ❖ Save and restore registers across procedure calls
 - ❖ Implement recursion
- ❖ Runtime stack is implemented via software convention
 - ❖ The **stack pointer** $\$sp = \29 (points to top of stack)
 - ❖ The **frame pointer** $\$fp = \30 (points to a procedure frame)

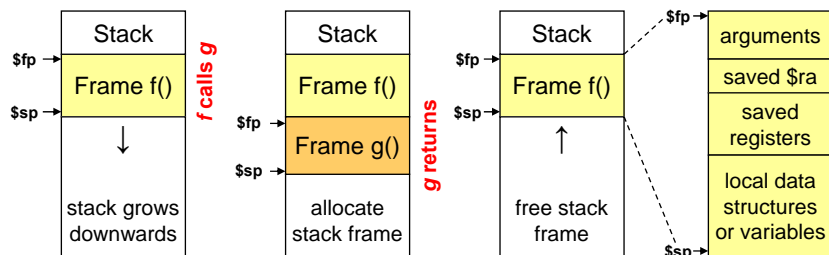
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Stack Frame

- ❖ **Stack frame** is the segment of the stack containing ...
 - ❖ Saved arguments, registers, and local data structures (if any)
- ❖ Called also the **activation frame** or **activation record**
- ❖ Frames are pushed and popped by adjusting ...
 - ❖ Stack pointer $\$sp = \29 and Frame pointer $\$fp = R30$
 - ❖ Decrement $\$sp$ to allocate stack frame, and increment to free



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Preserving Registers

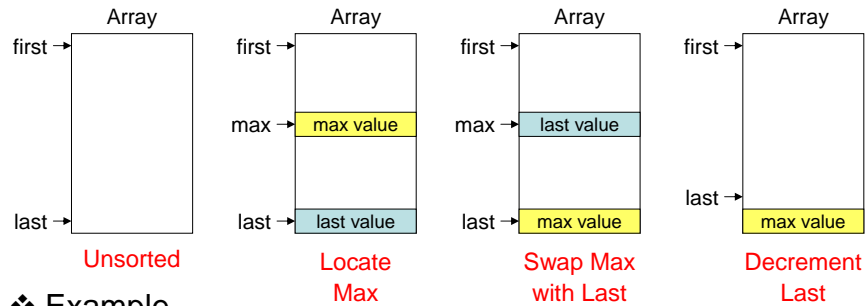
- ❖ Need to preserve registers across a procedure call
 - ❖ Stack can be used to preserve register values
- ❖ Which registers should be saved?
 - ❖ Registers modified by the called procedure, and
 - ❖ Still used by the calling procedure
- ❖ Who should preserve the registers?
 - ❖ **Called Procedure:** preferred method for modular code
 - Register preservation is done inside the called procedure
 - ❖ By convention, registers $\$s0$, $\$s1$, ..., $\$s7$ should be preserved
 - ❖ Also, registers $\$sp$, $\$fp$, and $\$ra$ should also be preserved

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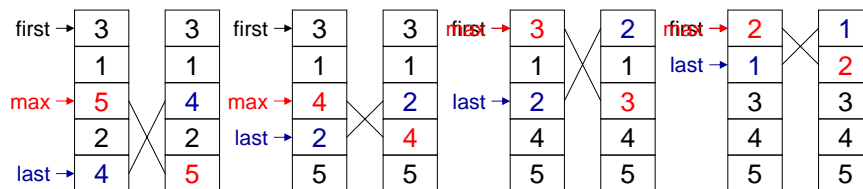
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Selection Sort



❖ Example



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Selection Sort Procedure

```

# Objective: Sort array using selection sort algorithm
#   Input: $a0 = pointer to first, $a1 = pointer to last
#   Output: array is sorted in place
#####
sort: addiu $sp, $sp, -4      # allocate one word on stack
      sw     $ra, 0($sp)     # save return address on stack
top:  jal    max             # call max procedure
      lw     $t0, 0($a1)     # $t0 = last value
      sw     $t0, 0($v0)    # swap last and max values
      sw     $v1, 0($a1)
      addiu $a1, $a1, -4    # decrement pointer to last
      bne   $a0, $a1, top   # more elements to sort
      lw     $ra, 0($sp)    # pop return address
      addiu $sp, $sp, 4
      jr    $ra             # return to caller
    
```

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Max Procedure

```
# Objective: Find the address and value of maximum element
#   Input: $a0 = pointer to first, $a1 = pointer to last
#   Output: $v0 = pointer to max, $v1 = value of max
#####
max:  move  $v0, $a0      # max pointer = first pointer
      lw   $v1, 0($v0)   # $v1 = first value
      beq  $a0, $a1, ret # if (first == last) return
      move $t0, $a0      # $t0 = array pointer
loop: addi $t0, $t0, 4   # point to next array element
      lw   $t1, 0($t0)   # $t1 = value of A[i]
      ble  $t1, $v1, skip # if (A[i] <= max) then skip
      move $v0, $t0      # found new maximum
      move $v1, $t1
skip: bne  $t0, $a1, loop # loop back if more elements
ret:  jr   $ra
```

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Example of a Recursive Procedure

```
int fact(int n) { if (n<2) return 1; else return (n*fact(n-1)); }
```

```
fact: slti  $t0,$a0,2    # (n<2)?
      beq  $t0,$0,else   # if false branch to else
      li   $v0,1         # $v0 = 1
      jr   $ra          # return to caller
else: addiu $sp,$sp,-8   # allocate 2 words on stack
      sw   $a0,4($sp)    # save argument n
      sw   $ra,0($sp)    # save return address
      addiu $a0,$a0,-1   # argument = n-1
      jal  fact          # call fact(n-1)
      lw   $a0,4($sp)    # restore argument
      lw   $ra,0($sp)    # restore return address
      mul  $v0,$a0,$v0   # $v0 = n*fact(n-1)
      addi $sp,$sp,8     # free stack frame
      jr   $ra          # return to caller
```

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