

LAB 07: MIPS Functions and Stack Segment

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Agenda

- 1 Caller vs. Callee
- 2 Functions
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- 5 Examples
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Function Example:

```
function: # function name  
# function body  
# return statement(essential!)  
jr $ra
```

Registers Convention

Register Name	Register No.	Register Usage
\$zero	\$0	Always zero, forced by hardware
\$at	\$1	Assembler Temporary register, reserved for assembler use
\$v0 - \$v1	\$2 - \$3	Results of a function
\$a0 - \$a3	\$4 - \$7	Arguments of a function
\$t0 - \$t7	\$8 - \$15	Registers for storing temporary values
\$s0 - \$s7	\$16 - \$23	Registers that should be saved across function calls
\$t8 - \$t9	\$24 - \$25	Registers for storing more temporary values
\$k0 - \$k1	\$26 - \$27	Registers reserved for the OS kernel use
\$gp	\$28	Global Pointer register that points to global data
\$sp	\$29	Stack Pointer register that points to top of stack
\$fp	\$30	Frame Pointer register that points to stack frame
\$ra	\$31	Return Address register used to return from a function call

Stack Segment

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Stack Segment

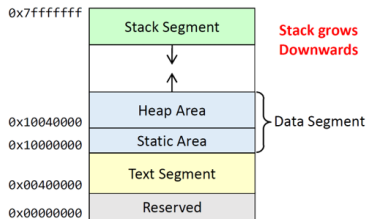
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MIPS Memory Organization

Recursive Function Example

```
int fact(int n){
    if (n < 2)
        return 1;
    else
        return n*fact(n-1);
}
```

Recursive factorial function

Recursive factorial function in MIPS Assembly

```
fact:
    bge $a0, 2, else    # branch if (n >= 2) to else
    li $v0, 1           # $v0 = 1
    jr $ra              # return to caller
else:
    addi $sp, $sp, -8    # allocate 8 bytes in the stack
    sw $a0, 0($sp)       # save the argument n
    sw $ra, 4($sp)       # save the return address
    addi $a0, $a0, -1    # argument $a0 = n-1
    jal fact             # call fact(n-1)
    lw $a0, 0($sp)       # restore $a0 = n
    lw $ra, 4($sp)       # restore return address
    mul $v0, $a0, $v0    # $v0 = n * fact(n-1)
    addi $sp, $sp, 8     # free stack frame
    jr $ra              # return to the caller
```

Live Examples

Task #1

Write a MIPS assembly program that implements the read, reverse, and print functions used by f function in Figure 7.6 & Figure 7.7 in the PDF file. These functions should work with any size n (not only size 10). Then write a main function that calls function f.

Example function	Stack Frame
<pre>void f() { int array[10]; read(array, 10); reverse(array, 10); print(array, 10); }</pre>	saved \$ra = 4 bytes
	<pre>int array[10] (40 bytes)</pre>

f function in C

```
f:   addiu $sp, $sp, -44    # allocate stack frame = 44 bytes
      sw   $ra, 40($sp)    # save $ra on the stack
      move $a0, $sp        # $a0 = address of array on the stack
      li   $a1, 10         # $a1 = 10
      jal  read            # call function read
      move $a0, $sp        # $a0 = address of array on the stack
      li   $a1, 10         # $a1 = 10
      jal  reverse         # call function reverse
      move $a0, $sp        # $a0 = address of array on the stack
      li   $a1, 10         # $a1 = 10
      jal  print           # call function print
      lw   $ra, 40($sp)    # load $ra from the stack
      addiu $sp, $sp, 44   # Free stack frame = 44 bytes
      jr   $ra             # return to caller
```

f function in MIPS Assembly

Task #1

Sample Run

```
Enter integer 1: 1
Enter integer 2: 2
Enter integer 3: 3
Enter integer 4: 4
Enter integer 5: 5
Enter integer 6: 6
Enter integer 7: 7
Enter integer 8: 8
Enter integer 9: 9
Enter integer 10: 10
Integer reversed = 10 9 8 7 6 5 4 3 2 1
```

Task #2

Write a MIPS assembly program that asks the user for an integer **n** he wishes to compute the Fibonacci number at that index. Calculate **fib(n)** based on the following code. Finally, print out the result.

```
int fib(int n) {  
    if (n <= 1)  
        return n;  
    return fib(n-1)+fib(n-2);  
}
```

Recursive Fibonacci function

Sample Run

```
Enter n: 7  
fib(n) = 13
```